

# Kinetic Backtracking of RNA folding

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## 1 Introduction

## 2 Algorithm

- Select set of states
- Find starting structure
- Extending the Front
- Saddle Height

## 3 Future Work

## Leipzig Group

### Chair of Bioinformatics

Konstantin Klemm: Networks

Matthias Kruspe: Sequence Alignments

Christoph Flamm: RNA Folding Dynamics, Energy Landscapes

### Chair of parallel Computing and Complex Systems

Michael Geis: RNA Folding, Swarm Intelligence

### Chair of Image Processing

Christian Heine: BarVis

Sebastian Poetzsch: PifoldVis

## Aim: Simulate the folding path of an RNA molecule

=> start at random coil

=> end at minimum free energy (mfE)

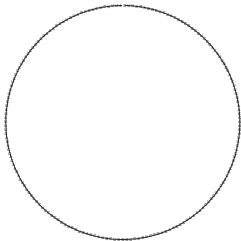
## Essential components:

**Select a set of states**

**Model state transitions**

Main issue: How to calculate energy barriers between states.

How do we get from A to B?



## Calculate the suboptimal structures

1. Find  $mfE$  via DP and backtrack suboptimal structures
2. This yields an upper triangular matrix with entry  $(i,j)$  denoting the free energy of the suboptimal structure on the subsequence  $(i,j)$
3. Filter out subsequences, in which base  $i$  and  $j$  pair.

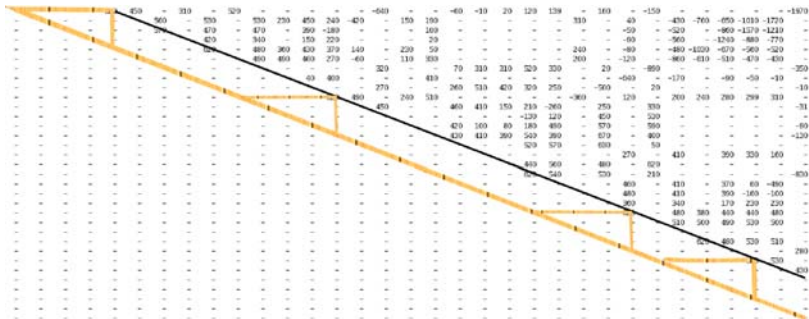




Table of integers, likely representing state transitions or energy values. The table is 100 rows by approximately 110 columns. The values range from -800 to 190. The data is sparse, with many empty cells.

## Determine the initial energy front

1. Rank extrema by
  - i) diagonal
  - ii) distance from edge of matrix
  - iii) 5' before 3'
2. Find the diagonal closest to the main diagonal that is not empty
3. Add substructures in that diagonal to the front if there is no conflict



## Conflict

A substructure  $(i,j)$  determines how the bases in that subsequence pair. Once  $(i,j)$  is part of the front, a pair  $(k,l)$  can only be added if  $[i, j] \cap [k, l] = \emptyset$

GUCCUUGCGUGAGGACAGCCCUUAUGUGAGGGC

.....  
 (...)..... ( 1, 5)  
 .....(...)..... (10,14)  
 .....(...)..... (22,26)  
 .....(...).. (27,31)  
 (...)....(...).....(...).....

## Algorithm

*saddleheight* = 0

*Front*  $F$  = *initial structure*

*while*( $F \neq mfE$ )

    increment *saddleheight*

    for  $x$  in *extrema*

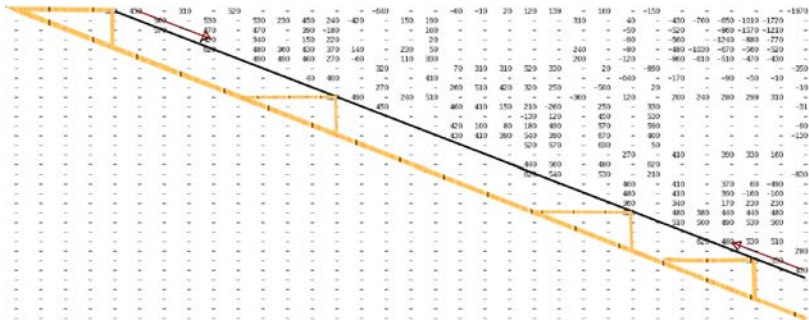
        calculate saddle point  $S$  between  $x$  and  $F$

*if*( $E(S) - E(F) < \textit{saddleheight}$ )

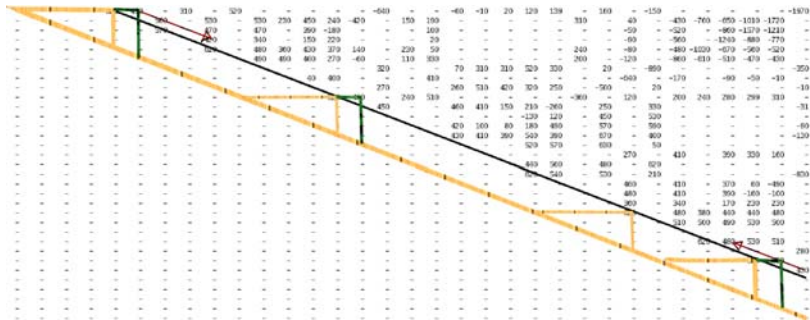
            extend  $F$  by  $x$

        remove all  $x$  in *extrema* that conflict with  $F$

# Extend Front

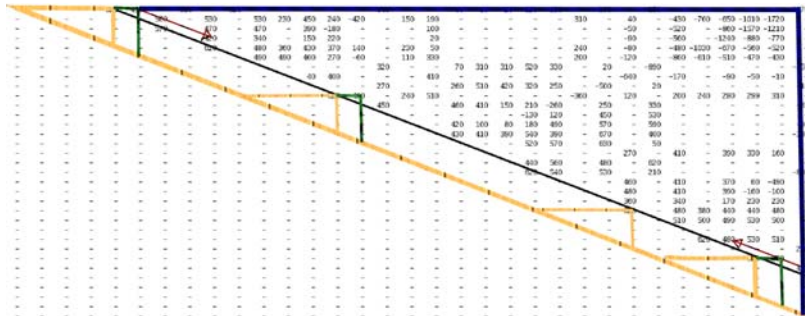


# Extend Front





# Final Front



## Saddle Height

To go from conformation A to B, there are  $(A\Delta B)!$  direct paths.

The best path does not have to be direct.

The saddle height is the highest energy of the lowest path:

$$S = \min_{P: A \rightarrow B} \max_{x \in P} E(x)$$

## Current Heuristic: Morgan-Higgs(1998)

1. Rank the elements in  $B \setminus A$  by conflict
2. For each  $x$  in  $B \setminus A$ 
  - i) Remove the base pairs in  $A$  that conflict with  $x$
  - ii) Add  $x$
  - iii) Add all other elements in  $B \setminus A$  that can be added now
  - iv) Record the energies of the traversed states
3. Take the highest recorded energy as saddle energy

## Future Work

Experiment with different heuristics for saddle heights

In Particular:

What is the lowest saddle point for  $n$  extensions of the front?

# Acknowledgments

Christoph Flamm, Peter Stadler